

The Ott L^AT_EX Layout Package `ottlayout.sty`

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November 29, 2019

1 Introduction

The Ott L^AT_EX Layout Package, `ottlayout.sty`, provides a range of options to tune the typesetting of Ott-generated inductive definition rules and grammars, overriding the default typesetting of the Ott-generated L^AT_EX code.

This document illustrates the common-case usage of the package, using Lightweight Java (LJ) [1] as an example Ott project. It should be read in conjunction with the source for this document (`manual.tex`) and the `Makefile`.

2 Usage

To use the package, one first uses Ott to generate L^AT_EX code with a chosen `-tex_name_prefix`, by default `ott`, but here `lj`, as in the example `Makefile`:

```
lj_included.tex : $(lj)
    ott $(INC_ARGS) -tex_name_prefix lj -tex $@ \
    -merge true $(lj)
```

Then one builds a file such as `lj_override.tex`, e.g. as below.

```
%_override.tex: override.tex empty.ott
    ott $(INC_ARGS) -tex_name_prefix lj \
    -tex_filter override.tex $@ empty.ott
```

This file simply contains redefinitions of some default L^AT_EX commands generated by Ott (with the `lj` prefix) to use the `ottlayout.sty` commands, e.g.

```
\renewcommand{\ljpremise}[1]{\premiseSTY{#1}}
\renewcommand{\ljusedrule}[1]{\usedruleSTY{#1}}
\renewcommand{\ljdrule}[4] [] {\druleSTY{#1}{#2}{#3}{#4}}
\renewenvironment{ljdefnblock}[3] [] {%
    \defnblockSTY{#1}{#2}{#3}}{\enddefnblockSTY}
```

Finally, in the user L^AT_EX document (for example this `manual.tex`), one: (a) includes the generated L^AT_EX for the user language, e.g. with `\include{lj_included}`; (b) uses the `ottlayout.sty` package, e.g. with `\usepackage{ottlayout}`; and (c) uses the generated override file to link the generated L^AT_EX with `ottlayout.sty`, e.g. with `\include{lj_override}`.

3 Displaying grammar

To display all Ott-generated L^AT_EX for LJ, we would normally write the command `\ljall{}`. To output all the LJ's grammar, we would use the L^AT_EX command: `\ljgrammar{}`. To show only selected parts of the grammar, we would normally use the command `\ljgrammartabular{}`. For example, to display the grammar of LJ's statement (`s`) and class definition (`cld`), we would write ¹

```
\ljgrammartabular{\ljs\ljinterrule\ljcld\ljafterlastrule}
```

to produce:

<i>s</i>	<code>::=</code>	statement
	<code>{ $\overline{s_k}^k$ }</code>	block
	<code>var = x;</code>	variable assignment
	<code>var = x.f;</code>	field read
	<code>x.f = y;</code>	field write
	<code>if (x == y) s else s'</code>	conditional branch
	<code>var = x.meth(\overline{y});</code>	method call
	<code>var = new_{ctx} cl();</code>	object creation
<i>cld</i>	<code>::=</code>	class
	<code>class dcl extends cl{$\overline{fd\ meth_def}$}</code>	def.

Alternatively, we can use the `ottlayout` package's `\grammartabularSTY{}` to produce a slightly more compact output, usually more suitable for publications. To display the same grammars, we would write

```
\grammartabularSTY{\ljs\\ \ljcld}
```

to produce:

<i>s</i>	<code>::=</code>	statement
	<code>{ $\overline{s_k}^k$ }</code>	block
	<code>var = x;</code>	variable assignment
	<code>var = x.f;</code>	field read
	<code>x.f = y;</code>	field write
	<code>if (x == y) s else s'</code>	conditional branch
	<code>var = x.meth(\overline{y});</code>	method call
	<code>var = new_{ctx} cl();</code>	object creation
<i>cld</i>	<code>::=</code>	class
	<code>class dcl extends cl{$\overline{fd\ meth_def}$}</code>	def.

¹The automatically generated grammar tabular command, here `\ljgrammartabular{}`, uses the `supertabular` package. Therefore, if we use the default grammar tabular, we have to explicitly import this package by writing `\usepackage{supertabular}` in our L^AT_EX document's prelude.

Note that in both cases the comments on the right are aligned according to the longest production in the block. Therefore, if the length of productions varies a lot, it is sometimes suitable to split them up into separate grammar tabulars. We could split the above example by writing

```
\grammartabularSTY{\ljs}\l
\grammartabularSTY{\ljcld}
```

to produce:

```
s ::=                                statement
|   {  $\overline{s_k}^k$  }           block
|   var = x;                       variable assignment
|   var = x.f;                     field read
|   x.f = y;                       field write
|   if (x == y) s else s'        conditional branch
|   var = x.meth( $\overline{y}$ );             method call
|   var = newctx cl();           object creation

cld ::=                             class
|   class dcl extends cl{ $\overline{fd\ meth\_def}$ }  def.
```

4 Displaying rules

To display all the rules of LJ, we would normally use the Ott-generated \LaTeX command `\ljdefnss{}`. To show all the rules of the LJ's reduction relation, we would use the command `\ljdefnrXXstmt{}` — if you are not sure what the name of the command you are looking for is, the easiest way to find out is to check the Ott-generated \LaTeX file, which is in our case `lj_included.tex`.

The `ottlayout` package gives many different options for displaying a particular rule and groups of rules. The currently available display options are:

Setting name	Possible values	Default value
showruleschema	yes no	yes
showcomment	yes no	yes
rulelayout	oneperline nobreaks	oneperline
premiselayout	oneperline oneline justify	justify
premisenamelayout	right left topright none	right
numberpremises	yes no	no
numbercolour	<i>any dvips colour name</i>	Gray

The default settings result in the same output as if the `ottlayout` package was not used.

We use LJ's (fairly complicated) reduction rule for methods to demonstrate a few of the available display settings for an example. To display the LJ rule `r_mcall` with default settings we write

\ljdrulerXXmcall{}

which produces:

$$\begin{array}{c}
L(x) = oid \\
H(oid) = \tau \\
\textbf{find_meth_def} (P, \tau, meth) = (ctx, cl \textit{meth}(\overline{cl_k \textit{var}_k}^k) \{ \overline{s_j'}^j \textbf{return } y; \}) \\
\overline{var_k'}^k \perp \textbf{dom} (L) \\
\textbf{distinct} (\overline{var_k'}^k) \\
x' \notin \textbf{dom} (L) \\
x' \notin \overline{var_k'}^k \\
\overline{L(y_k) = v_k}^k \\
L' = L[\overline{var_k' \mapsto v_k}^k][x' \mapsto oid] \\
\theta = [\overline{var_k \mapsto var_k'}^k][\textbf{this} \mapsto x'] \\
\overline{\theta \vdash s_j' \rightsquigarrow s_j''^j} \\
\theta(y) = y'
\end{array}
\frac{}{(P, L, H, var = x.meth(\overline{y_k}^k); \overline{s_l^l}) \longrightarrow (P, L', H, \overline{s_j''^j} \textit{var} = y'; \overline{s_l^l})} \text{R_MCALL}$$

Note that math mode is entered automatically, which means that we can use any L^AT_EX text layout utilities to layout our rules as we wish.

We can change the default setting with command \ottstyledefaults{} by passing keys and values in the KeyVal style. For example, to make the premises display more compactly in all rules from now on, we write

\ottstyledefaults{premiselayout=justify}

Now the the command \ljdruleXXmcall{} produces the following instead:

$$\begin{array}{c}
L(x) = oid \qquad H(oid) = \tau \\
\textbf{find_meth_def} (P, \tau, meth) = (ctx, cl \textit{meth}(\overline{cl_k \textit{var}_k}^k) \{ \overline{s_j'}^j \textbf{return } y; \}) \\
\overline{var_k'}^k \perp \textbf{dom} (L) \quad \textbf{distinct} (\overline{var_k'}^k) \quad x' \notin \textbf{dom} (L) \\
x' \notin \overline{var_k'}^k \quad \overline{L(y_k) = v_k}^k \\
L' = L[\overline{var_k' \mapsto v_k}^k][x' \mapsto oid] \quad \theta = [\overline{var_k \mapsto var_k'}^k][\textbf{this} \mapsto x'] \\
\overline{\theta \vdash s_j' \rightsquigarrow s_j''^j} \qquad \theta(y) = y'
\end{array}
\frac{}{(P, L, H, var = x.meth(\overline{y_k}^k); \overline{s_l^l}) \longrightarrow (P, L', H, \overline{s_j''^j} \textit{var} = y'; \overline{s_l^l})} \text{R_MCALL}$$

To use the non-default settings for a particular rule, we can write the same KeyVal pairs inside as the parameter to the rule command. For example, to number the premises in the rule, to make the numbers yellow-orange, and to place the rule's name in top-right corner, we write

\ljdrulerXXmcall{numberpremises=yes,
numbercolour=YellowOrange,
premisenamelayou=topright}

which produces:

R_MCALL

$$\begin{array}{c}
\begin{array}{ll}
1. L(x) = oid & 2. H(oid) = \tau \\
3. \text{find_meth_def } (P, \tau, meth) = (ctx, cl \text{ meth } (\overline{cl_k \text{ var}_k})^k) \{ \overline{s_j'}^j \text{ return } y; \} \\
4. \overline{var_k'}^k \perp \text{dom}(L) & 5. \text{distinct } (\overline{var_k'}^k) \quad 6. x' \notin \text{dom}(L) \\
7. x' \notin \overline{var_k'}^k & 8. \overline{L(y_k) = v_k}^k \\
9. L' = L[\overline{var_k'} \mapsto v_k]^k [x' \mapsto oid] \\
10. \theta = [\overline{var_k} \mapsto \overline{var_k'}^k] [\text{this} \mapsto x'] & 11. \overline{\theta \vdash s_j' \rightsquigarrow s_j''^j} \\
12. \theta(y) = y'
\end{array} \\
\hline
(P, L, H, var = x.meth(\overline{y_k}^k); \overline{s_l}^l) \longrightarrow (P, L', H, \overline{s_j''}^j \text{ var} = y'; \overline{s_l}^l)
\end{array}$$

As you can see, the setting for compacting the premises was kept, because it was set globally for all rules following the `\ottstyledefaults{}` command.

As with commands for individual rules, we can pass in KeyVal pairs to the commands that display groups of rules, which will affect how the rules of that particular group of rules is displayed. If we wanted to display LJ's reduction rules for statements with current default display settings, we would write `\ljdefnrXXstmt{}`. To display the reduction rules with rule names on the left side, and their premises numbered, we write

`\ljdefnrXXstmt{numberpremises=yes, premisenamelayou=left}`

This produces²:

$$\boxed{config \longrightarrow config'} \quad \text{one step reduction}$$

$$\begin{array}{l}
\text{R_FIELD_READ_NPE} \quad \frac{1. L(x) = \mathbf{null}}{(P, L, H, var = x.f; \overline{s_l}^l) \longrightarrow (P, L, H, \mathbf{NPE})} \\
\text{R_FIELD_WRITE_NPE} \quad \frac{1. L(x) = \mathbf{null}}{(P, L, H, x.f = y; \overline{s_l}^l) \longrightarrow (P, L, H, \mathbf{NPE})} \\
\text{R_VAR_ASSIGN} \quad \frac{1. L(x) = v}{(P, L, H, var = x; \overline{s_l}^l) \longrightarrow (P, L[var \mapsto v], H, \overline{s_l}^l)} \\
\text{R_FIELD_READ} \quad \frac{1. L(x) = oid \quad 2. H(oid, f) = v}{(P, L, H, var = x.f; \overline{s_l}^l) \longrightarrow (P, L[var \mapsto v], H, \overline{s_l}^l)} \\
\text{R_FIELD_WRITE} \quad \frac{1. L(x) = oid \quad 2. L(y) = v}{(P, L, H, x.f = y; \overline{s_l}^l) \longrightarrow (P, L, H[(oid, f) \mapsto v], \overline{s_l}^l)} \\
\text{R_MCALL_NPE} \quad \frac{1. L(x) = \mathbf{null}}{(P, L, H, var = x.meth(\overline{y_k}^k); \overline{s_l}^l) \longrightarrow (P, L, H, \mathbf{NPE})}
\end{array}$$

²We set the font size to `small` so that the rules are not too wide.

$$\begin{array}{l}
\text{R_MCALL_CNFE} \quad \frac{
\begin{array}{l}
1. L(x) = oid \quad 2. H(oid) = \tau \\
3. \mathbf{find_meth_def} (P, \tau, meth) = \mathbf{null}
\end{array}
}{(P, L, H, var = x.meth(\overline{y_k^k}); \overline{s_l^l}) \longrightarrow (P, L, H, \mathbf{CNFE})} \\
\\
\text{R_IF_TRUE} \quad \frac{
\begin{array}{l}
1. L(x) = v \quad 2. L(y) = w \quad 3. v == w
\end{array}
}{(P, L, H, \mathbf{if} (x == y) s_1 \mathbf{else} s_2 \overline{s_l^l}) \longrightarrow (P, L, H, s_1 \overline{s_l^l})} \\
\\
\text{R_IF_FALSE} \quad \frac{
\begin{array}{l}
1. L(x) = v \quad 2. L(y) = w \quad 3. v \neq w
\end{array}
}{(P, L, H, \mathbf{if} (x == y) s_1 \mathbf{else} s_2 \overline{s_l^l}) \longrightarrow (P, L, H, s_2 \overline{s_l^l})} \\
\\
\text{R_BLOCK} \quad \frac{}{(P, L, H, \{ \overline{s_k^k} \} \overline{s_l^l}) \longrightarrow (P, L, H, \overline{s_k^k} \overline{s_l^l})} \\
\\
\text{R_NEW} \quad \frac{
\begin{array}{l}
1. \mathbf{find_type} (P, ctx, cl) = \tau \\
2. \mathbf{fields} (P, \tau) = \overline{f_k^k} \quad 3. oid \notin \mathbf{dom} (H) \\
4. H' = H[oid \mapsto (\tau, \overline{f_k} \mapsto \mathbf{null}^k)]
\end{array}
}{(P, L, H, var = \mathbf{new}_{ctx} cl(); \overline{s_l^l}) \longrightarrow (P, L[var \mapsto oid], H', \overline{s_l^l})} \\
\\
\text{R_MCALL} \quad \frac{
\begin{array}{l}
1. L(x) = oid \quad 2. H(oid) = \tau \\
3. \mathbf{find_meth_def} (P, \tau, meth) = (ctx, cl \text{ meth}(\overline{cl_k} \overline{var_k^k}) \{ \overline{s_j^{t_j}} \mathbf{return} y; \}) \\
4. \overline{var_k^{t_k}} \perp \mathbf{dom} (L) \quad 5. \mathbf{distinct} (\overline{var_k^{t_k}}) \quad 6. x' \notin \mathbf{dom} (L) \\
7. x' \notin \overline{var_k^{t_k}} \quad 8. \overline{L(y_k)} = v_k \\
9. L' = L[\overline{var_k^{t_k}} \mapsto v_k][x' \mapsto oid] \\
10. \theta = [\overline{var_k} \mapsto \overline{var_k^{t_k}}][\mathbf{this} \mapsto x'] \quad 11. \theta \vdash s_j' \rightsquigarrow s_j''^{t_j} \\
12. \theta(y) = y'
\end{array}
}{(P, L, H, var = x.meth(\overline{y_k^k}); \overline{s_l^l}) \longrightarrow (P, L', H, \overline{s_j^{t_j}} var = y'; \overline{s_l^l})}
\end{array}$$

Therefore, Ott-generated \LaTeX commands for rules of a specific **defn** in the Ott source file can take KeyVal arguments. However, note that currently the Ott-generated \LaTeX output does not allow for this package to allow the same arguments being passed to commands for a group of **defns**, i.e. **defns**.

References

- [1] STRNIŠA, R., AND PARKINSON, M. Lightweight Java. <http://www.cl.cam.ac.uk/~rs456/1j>, Sept. 2006.